

Stability Analysis of Networked Control Systems

A Sum of Squares Approach

Nicolas Bauer Paul Maas Maurice Heemels

Hybrid & Networked Systems Group
Department of Mechanical Engineering

CDC, Atlanta - WeC19.2
December 15, 2010

TU / **e**

Technische Universiteit
Eindhoven
University of Technology

Where innovation starts

Network effects:

- ▶ varying transmission intervals $h_k \in [h_{min}, h_{max}], \forall k$
- ▶ varying delays $\tau_k \in [\tau_{min}, \tau_{max}]$, where $\tau_k \leq \min\{\tau_{max}, h_k\}, \forall k$
- ▶ communication constraints

→ Stability Analysis

COMMUNICATION CONSTRAINTS

- Node:** $\sigma \in \{1, \dots, N\}$ a collection of sensors and/or actuators that can transmit data simultaneously
- Protocol:** $\delta(\sigma, \mathbf{e})$ a function that determines which node is allowed to transmit data
- Network Error:** $\mathbf{e} \in \mathbb{R}^{n_e}$ $\mathbf{e} = [\mathbf{e}_u^\top, \mathbf{e}_y^\top]^\top$ where $\mathbf{e}_u = \hat{\mathbf{u}} - \mathbf{u}$ and $\mathbf{e}_y = \hat{\mathbf{y}} - \mathbf{y}$

RR protocol:

$$\sigma_{NEW} = \delta_{RR}(\sigma_{OLD}, \mathbf{e}) := (\sigma \bmod N) + 1$$

TOD/MEF protocol:

$$\sigma_{NEW} = \delta_{TOD}(\sigma_{OLD}, \mathbf{e}) := \arg \max_j |e_j|, \quad j = 1, \dots, N$$

At updates at times $t_k + \tau_k$, $k \in \mathbb{N}$

$$\hat{y}((t_k + \tau_k)^+) = \mathbf{y}(t_k) + \tilde{\mathbf{h}}(\delta(\sigma(t_k), \mathbf{e}(t_k)), \mathbf{e}(t_k))$$

where

$$\tilde{\mathbf{h}}_i(\delta(\sigma(t_k), \mathbf{e}(t_k)), \mathbf{e}(t_k)) := \begin{cases} 0, & \text{if } i = \delta(\sigma(t_k), \mathbf{e}(t_k)), \\ \mathbf{e}_i(t_k), & \text{otherwise} \end{cases}$$



$$\begin{aligned}\dot{\xi} &= F(\xi), \quad \xi \in C, \\ \xi^+ &= G(\xi), \quad \xi \in D\end{aligned}$$

[1] Goebel, Sanfelice, and Teel. *IEEE Control Systems Magazine*, April 2009.

[2] Nešić and Teel. *IEEE Transactions on Automatic Control*, October 2004.

[3] Heemels et al. *IEEE Transactions on Automatic Control*, August 2010.

State variables:

- ▶ $x \in \mathbb{R}^{n_x}$ - combined state $[x_p^\top, x_c^\top]^\top$ of the plant and controller
- ▶ $e \in \mathbb{R}^{n_e}$ - the network-induced error $[e_u^\top, e_y^\top]^\top$
- ▶ $s \in \mathbb{R}^{n_e}$ - memory storing the value $h(\sigma, e) - e$ at t_k
- ▶ $\tau \in \mathbb{R}_{\geq 0}$ - timer to guarantee $\tau_k \in [\tau_{min}, \tau_{max}]$, $h_k \in [h_{min}, h_{max}]$
- ▶ $\sigma \in \{1, \dots, N\}$ - latest node that got access to the network
- ▶ $\ell \in \{0, 1\}$ - signifies a transmission or an update

the entire state: $\xi = [x^\top, e^\top, s^\top, \tau, \sigma, \ell]^\top \in \mathbb{R}^{n_\xi}$

The hybrid system model

$$\begin{aligned}\dot{\xi} &= F(\xi), \quad \xi \in C, \\ \xi^+ &= G(\xi), \quad \xi \in D\end{aligned}$$

where

$$C := \{\xi \in \mathbb{R}^{n_\xi} \mid (\ell = 0 \wedge \tau \in [0, h_{max}]) \vee (\ell = 1 \wedge \tau \in [0, \tau_{max}])\}$$

$$F(\xi) = F(x, e, s, \tau, \sigma, \ell) := (f(x, e), g(x, e), 0, 1, 0, 0)$$

and

$$D := \{\xi \in \mathbb{R}^{n_\xi} \mid (\ell = 0 \wedge \tau \in [h_{min}, h_{max}]) \vee (\ell = 1 \wedge \tau \in [\tau_{min}, \tau_{max}])\}$$

$$G(\xi) := \begin{cases} G(x, e, s, \tau, \sigma, 0) &= (x, e, \tilde{h}(\delta(\sigma, e), e) - e, 0, \delta(\sigma, e), 1) & \{t_k\} \\ G(x, e, s, \tau, \sigma, 1) &= (x, s - e, 0, \tau, \sigma, 0) & \{t_k + \tau_k\} \end{cases}$$

THEOREM: Consider a hybrid system $\Sigma = (C, F, D, G)$ and a compact set $\mathcal{A} \subset \mathbb{R}^{n_\xi}$ satisfying $G(D \cap \mathcal{A}) \subset \mathcal{A}$. If every solution of Σ exists for all times $t \in [0, \infty)$ and there exists a Lyapunov function candidate V for (Σ, \mathcal{A}) that satisfies

$$\langle \nabla V(\xi), F(\xi) \rangle < 0 \text{ for all } \xi \in C \setminus \mathcal{A}$$

$$V(G(\xi)) - V(\xi) \leq 0 \text{ for all } \xi \in D \setminus \mathcal{A},$$

then the set \mathcal{A} is GAS.

THEOREM: Consider a hybrid system $\Sigma = (C, F, D, G)$ and a compact set $\mathcal{A} \subset \mathbb{R}^{n_\xi}$ satisfying $G(D \cap \mathcal{A}) \subset \mathcal{A}$. If every solution of Σ exists for all times $t \in [0, \infty)$ and there exists a Lyapunov function candidate V for (Σ, \mathcal{A}) that satisfies

$$\langle \nabla V(\xi), F(\xi) \rangle < 0 \text{ for all } \xi \in C \setminus \mathcal{A}$$

$$V(G(\xi)) - V(\xi) \leq 0 \text{ for all } \xi \in D \setminus \mathcal{A},$$

then the set \mathcal{A} is GAS.

DEFINITION: A polynomial $p(x)$ is called a sum of squares (SOS) if there exist polynomials $p_1(x), p_2(x), \dots, p_m(x)$ such that $p(x) = \sum_{i=1}^m p_i^2(x)$ for all $x \in \mathbb{R}^n$.

- ▶ SOS can be checked via LMI conditions!!

$$\begin{aligned}\dot{\xi} &= F(\xi), & \xi \in C, \\ \xi^+ &= G(\xi), & \xi \in D\end{aligned}$$



$$\begin{aligned}\dot{\xi} &= F_i(\xi), & \xi \in C_i, \\ \xi^+ &= G_m(\xi), & \xi \in D_m\end{aligned}$$

$$\begin{aligned}C_i &= \{\xi \in \mathbb{R}^{n_\xi} \mid c_{i,j}(\xi) \geq 0, \text{ for } j = 1, \dots, m_C^i, \\ &\quad \bar{c}_{i,l}(\xi) = 0, \text{ for } l = 1, \dots, n_C^i\}, \\ D_m &= \{\xi \in \mathbb{R}^{n_\xi} \mid d_{m,j}(\xi) \geq 0, \text{ for } j = 1, \dots, m_D^m, \\ &\quad \bar{d}_{m,l}(\xi) = 0, \text{ for } l = 1, \dots, n_D^m\}\end{aligned}$$

- ▶ basic semialgebraic sets

- ▶ Is $F(\xi)$ (piecewise) polynomial?

$$F(\xi) = F(x, e, s, \tau, \sigma, \ell) := (f(x, e), g(x, e), 0, 1, 0, 0)$$

$$C := \{\xi \in \mathbb{R}^{n_\xi} \mid (\ell = 0 \wedge \tau \in [0, h_{max}]) \vee (\ell = 1 \wedge \tau \in [0, \tau_{max}])\}$$

- ▶ Yes! ✓

- ▶ Is $F(\xi)$ (piecewise) polynomial?

$$F(\xi) = F(x, e, s, \tau, \sigma, \ell) := (f(x, e), g(x, e), 0, 1, 0, 0)$$

$$C := \{\xi \in \mathbb{R}^{n_\xi} \mid (\ell = 0 \wedge \tau \in [0, h_{max}]) \vee (\ell = 1 \wedge \tau \in [0, \tau_{max}])\}$$

- ▶ Yes! ✓

- ▶ Is $G(\xi)$ (piecewise) polynomial?

$$G(\xi) := \begin{cases} G(x, e, s, \tau, \sigma, 0) & = (x, e, \tilde{h}(\delta(\sigma, e), e) - e, 0, \delta(\sigma, e), 1) \\ G(x, e, s, \tau, \sigma, 1) & = (x, s - e, 0, \tau, \sigma, 0) \end{cases}$$

$$D := \{\xi \in \mathbb{R}^{n_\xi} \mid (\ell = 0 \wedge \tau \in [h_{min}, h_{max}]) \vee (\ell = 1 \wedge \tau \in [\tau_{min}, \tau_{max}])\}$$

- ▶ Doesn't appear so ...

IDEA: Encode the protocol $\delta(\sigma, e)$ into the jump set D

$G(\xi)$ is piecewise polynomial ...

$$G(x, e, s, \tau, j, 0) = (x, e, \tilde{h}(j, e) - e, 0, j, 1), \quad j = 1, \dots, N$$

$$G(x, e, s, \tau, \sigma, 1) = (x, s + e, 0, \tau, \sigma, 0)$$

for $j = 1, \dots, N$, corresponding jump map

IDEA: Encode the protocol $\delta(\sigma, e)$ into the jump set D

$G(\xi)$ is piecewise polynomial ...

$$G(x, e, s, \tau, j, 0) = (x, e, \tilde{h}(j, e) - e, 0, j, 1), \quad j = 1, \dots, N$$

$$G(x, e, s, \tau, \sigma, 1) = (x, s + e, 0, \tau, \sigma, 0)$$

for $j = 1, \dots, N$, corresponding jump map

$$RR \left\{ \begin{array}{l} D_{0,1} = \{ \xi \in \mathbb{R}^{n_\xi} \mid \ell = 0, \tau - h_{min} \geq 0, h_{max} - \tau \geq 0, \sigma = N \} \\ D_{0,j} = \{ \xi \in \mathbb{R}^{n_\xi} \mid \ell = 0, \tau - h_{min} \geq 0, h_{max} - \tau \geq 0, \sigma + 1 = j \} \end{array} \right.$$

IDEA: Encode the protocol $\delta(\sigma, e)$ into the jump set D

$G(\xi)$ is piecewise polynomial ...

$$G(x, e, s, \tau, j, 0) = (x, e, \tilde{h}(j, e) - e, 0, j, 1), \quad j = 1, \dots, N$$

$$G(x, e, s, \tau, \sigma, 1) = (x, s + e, 0, \tau, \sigma, 0)$$

for $j = 1, \dots, N$, corresponding jump map

$$\begin{array}{l}
 RR \\
 TOD
 \end{array}
 \left\{ \begin{array}{l}
 D_{0,1} = \{ \xi \in \mathbb{R}^{n_\xi} \mid \ell = 0, \tau - h_{min} \geq 0, h_{max} - \tau \geq 0, \sigma = N \} \\
 D_{0,j} = \{ \xi \in \mathbb{R}^{n_\xi} \mid \ell = 0, \tau - h_{min} \geq 0, h_{max} - \tau \geq 0, \sigma + 1 = j \} \\
 D_{0,j} = \{ \xi \in \mathbb{R}^{n_\xi} \mid \ell = 0, \tau - h_{min} \geq 0, h_{max} - \tau \geq 0, \\
 \qquad \qquad \qquad e_j^\top e_j - e_i^\top e_i \geq 0 \forall i = 1, \dots, N \}
 \end{array} \right.$$

IDEA: Encode the protocol $\delta(\sigma, e)$ into the jump set D

$G(\xi)$ is piecewise polynomial ...

$$G(x, e, s, \tau, j, 0) = (x, e, \tilde{h}(j, e) - e, 0, j, 1), \quad j = 1, \dots, N$$

$$G(x, e, s, \tau, \sigma, 1) = (x, s + e, 0, \tau, \sigma, 0)$$

for $j = 1, \dots, N$, corresponding jump map

$$\begin{array}{l} RR \\ TOD \end{array} \left\{ \begin{array}{l} D_{0,1} = \{ \xi \in \mathbb{R}^{n_\xi} \mid \ell = 0, \tau - h_{min} \geq 0, h_{max} - \tau \geq 0, \sigma = N \} \\ D_{0,j} = \{ \xi \in \mathbb{R}^{n_\xi} \mid \ell = 0, \tau - h_{min} \geq 0, h_{max} - \tau \geq 0, \sigma + 1 = j \} \\ D_{0,j} = \{ \xi \in \mathbb{R}^{n_\xi} \mid \ell = 0, \tau - h_{min} \geq 0, h_{max} - \tau \geq 0, \\ \quad \quad \quad e_j^\top e_j - e_i^\top e_i \geq 0 \forall i = 1, \dots, N \} \end{array} \right.$$

with

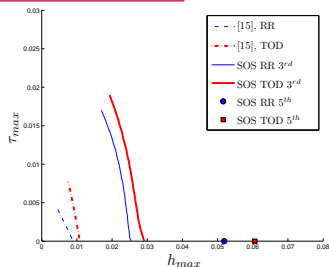
$$D_1 = \{ \xi \in \mathbb{R}^{n_\xi} \mid \ell = 1, \tau - \tau_{min} \geq 0, \tau_{max} - \tau \geq 0 \}$$

KEY STEPS:

1. NCS modeled as a hybrid system
2. LF candidate is structured polynomially \rightarrow Flexibility
3. SOS + Hybrid Stability Theorem \rightarrow LMI conditions for GAS
4. Embedded protocol in $D \rightarrow$ NCS is a piecewise poly hybrid system
5. regional info $C_i, D_m \rightarrow$ positivstellensatz (type of S-procedure)

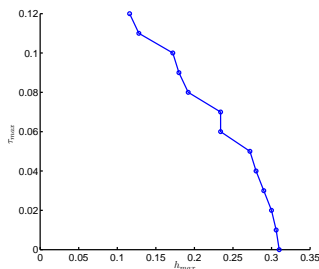
\rightarrow LMI conditions to prove GAS of the NCS!

BATCH REACTOR



- ▶ 4 state plant
- ▶ 2 state controller
- ▶ y transmitted over network
- ▶ u wired
- ▶ $\tau_{min} = h_{min} = 0$

POLYNOMIAL DYNAMICS



- $$\dot{x}(t) = -x^3(t) + x^2(t)u(t)$$
- $$u(t) = -x(t)$$
- ▶ no comm. constraints

Advantages

- ▶ automated method for stability analysis of NCSs
- ▶ less conservative than earlier methods
- ▶ piecewise polynomial plants and controllers
- ▶ non-zero lower bounds on varying delays and transmission intervals
- ▶ does not require an overapproximation of the NCS

Disadvantages

- ▶ computation complexity grows rapidly with state dimension
→ Improved solvers needed